Similarity Search
Trees

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Outline

1. What is a Tree?

2. Encoding XML as Trees

Outline

What is a Tree?

- Graph: a pair (N, E) of nodes N and edges E between nodes of N
- Tree: a finite, directed, acyclic graph T
  - that is connected and
  - no node has more than one incoming edge
- Edges: $E(T)$ are the edges of T
  - an edge $(p, c) \in E(T)$ is an ordered pair
  - with $p, c \in N(T)$
- "Special" Nodes: $N(T)$ are the nodes of T
  - parent/child: $(p, c) \in E(T) \Leftrightarrow p$ is the parent of c, c is the child of p
  - siblings: c_1 and c_2 are siblings if they have the same parent node
  - root node: node without parent (no incoming edge)
  - leaf node: node without children (no outgoing edge)
  - fanout: fanout f_v of node v is the number of children of v
  - descendants: node x is a descendant of v if x is reachable from v

Unlabeled Trees

- Unlabeled Tree:
  - the focus is on the structure, not on distinguishing nodes
  - however, we need to distinguish nodes in order to define edges
    $\Rightarrow$ each node v has a unique identifier id(v) within the tree

- Example: $T = (\{1, 3, 5, 4, 7\}, \{(1, 3), (1, 5), (5, 4), (5, 7)\})$
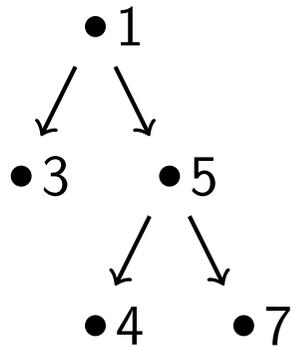
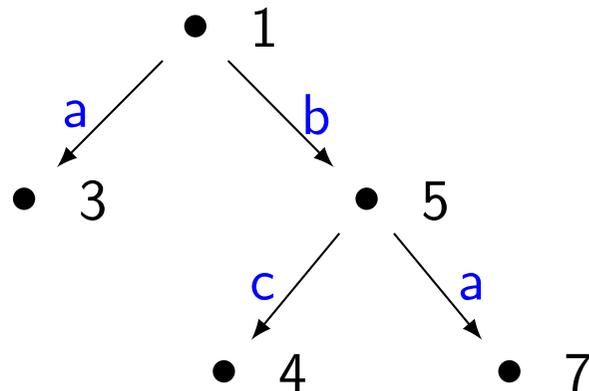# Edge Labeled Trees

- Edge Labeled Tree:
  - an edge $e \in E(T)$ between nodes a and b is a triple $e = (id(a), id(b), \lambda(e))$
  - $id(a)$ and $id(b)$ are node IDs
  - $\lambda(e)$ is the edge label (not necessarily unique within the tree)

- Example:

  $T = (\{1, 3, 5, 4, 7\}, \{(1, 3, a), (1, 5, b), (5, 4, c), (5, 7, a)\})$
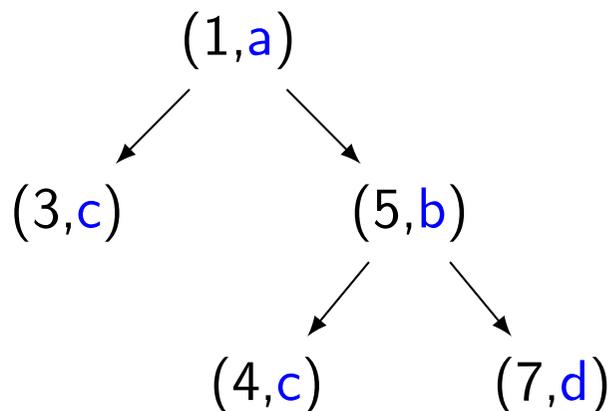
Node Labeled Trees

- Node Labeled Tree:
  - a node $v \in N(T)$ is a pair $(id(v), \lambda(v))$
  - $id(v)$ is unique within the tree
  - label $\lambda(v)$ needs not to be unique

- Intuition:
  - The identifier is the key of the node.
  - The label is the data carried by the node.
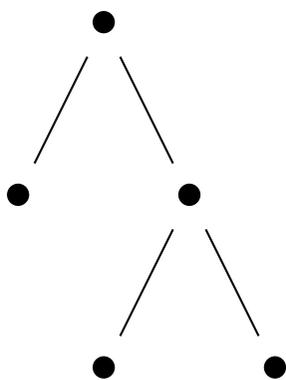
- Example: $T = (\{(1, a), (3, c), (5, b), (4, c), (7, d)\},$
  $\{(1, 3), (1, 5), (5, 4), (5, 7)\})$

(1,a)

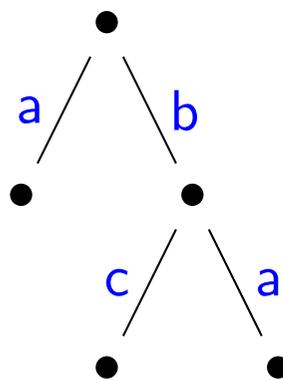(3,c)   (5,b)

(4,c)   (7,d)

Notation and Graphical Representation

- Notation:
  - node identifiers: $\text{id}(v_i) = i$
  - tree identifiers: $T_1, T_2, \ldots$

- Graphical representation
  - we omit brackets for (identifier,label)-pairs
  - we (sometimes) omit node identifiers at all
  - we do not show the direction of edges
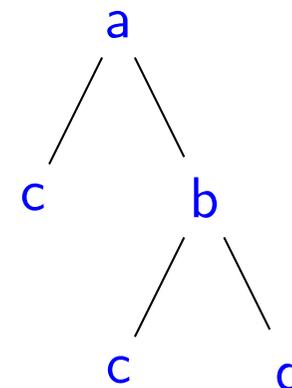   (edges are always directed from root to leave)
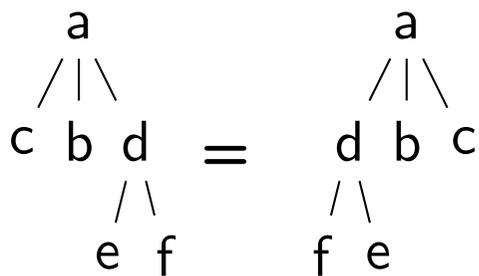
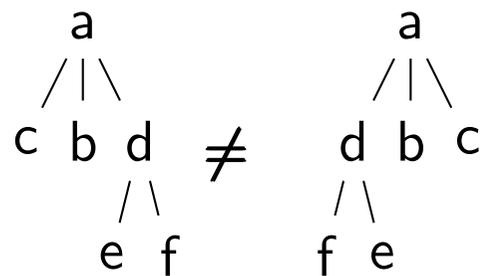| **unlabeled tree**	**edge labeled tree**	**node labeled tree**

Ordered Trees

- Ordered Trees: siblings are ordered
- Contiguous siblings $s_1 < s_2$ have no sibling x such that $s_1 < x < s_2$
- c_i is the *i*-th child of p if
  - p is the parent of c_i, and
  - $i = |\{x \in N(T) : (p, x) \in E(T), x \leq c_i\}|$
- Node order $<$ is a strict total order that respects the tree structure:
  - (a) If *c* is a child of *v*, then $c < v$
  - (b) If $f < g$ and *f* is not a descendant of *g*, then $f < g'$ for all descendants g' of g
- Example:

Unordered Trees	**Ordered Trees**

```
    Unordered Trees       |      Ordered Trees
   a            a         |     a            a
  /|\          /|\        |    /|\          /|\
 c b d  =   d b c         |   c b d  ≠   d b c
    /\        /\          |      /\        /\
   e f       f e          |     e f       f e
```

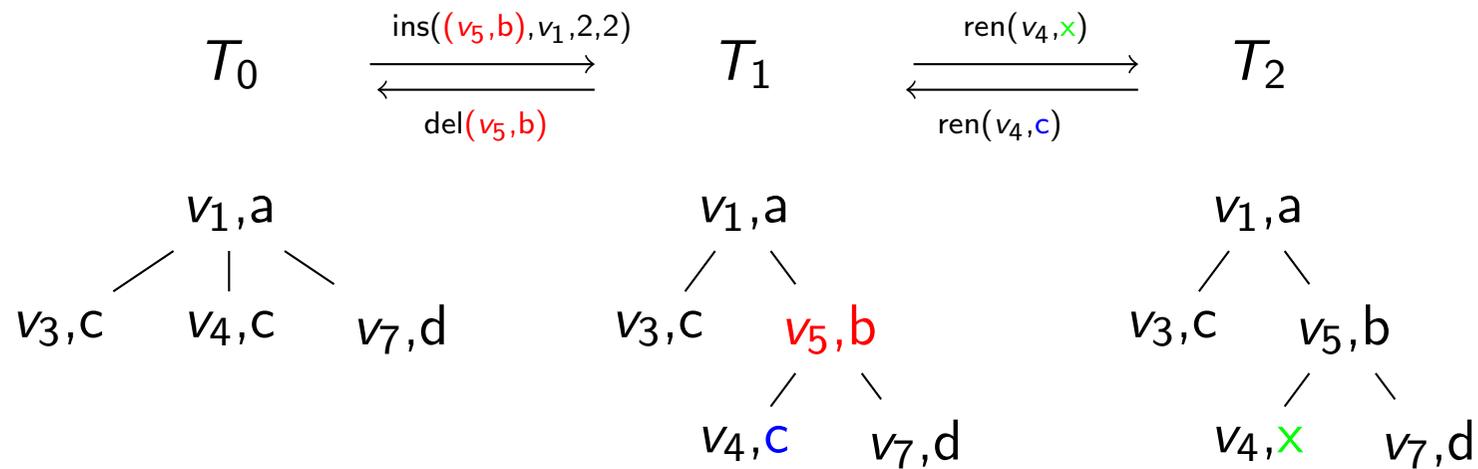- Note: "ordered" does not necessarily mean "sorted alphabetically"

Edit Operations

- We assume ordered, labeled trees
- Rename node: $ren(v, l')$
  - change label l of v to $l' \neq l$
- Delete node: $del(v)$ (v is not the root node)
  - remove v
  - connect v's children directly to v's parent node (preserving order)
- Insert node: $ins(v, p, k, m)$
  - detach m consecutive children of p, starting with the child at position k, i.e., the children $c_k, c_{k+1}, \ldots, c_{k+m-1}$
  - attach $c_k, c_{k+1}, \ldots, c_{k+m-1}$ as children of the new node v (preserving order)
  - insert new node v as k-th child of p
- Insert and delete are inverse edit operations (i.e., insert undoes delete and vice versa)

Example: Edit Operations

$$T_0 \xrightarrow[\text{del}(v_5,b)]{\text{ins}((v_5,b),v_1,2,2)} T_1 \xrightarrow[\text{ren}(v_4,c)]{\text{ren}(v_4,x)} T_2$$

T_0: v_1,a with children v_3,c — v_4,c — v_7,d

T_1: v_1,a with children v_3,c and v_5,b; v_5,b with children v_4,c and v_7,d

T_2: v_1,a with children v_3,c and v_5,b; v_5,b with children v_4,x and v_7,d

# Outline

Representing XML as a Tree

- Many possibilities – we will consider
  - single-label tree
  - double-label tree

- Pros/cons depend on application!

XML as a Single-Label Tree

- The XML document is encoded as a tree with:
  - XML element: node labeled with element tag name
  - XML attribute: node labeled with attribute name
  - Text contained in elements/attributes: node labeled with the text-value

- Element nodes contain:
  - nodes of their sub-elements
  - nodes of their attributes
  - nodes with their text values

- Attribute nodes contain:
  - single node with their text value

- Text nodes are always leaves

- Order:
  - sub-element and text nodes are ordered
  - attributes are not ordered (approach: store them before all sub-elements, sort according to attribute name)

Example: XML as a Single-Label Tree

```
<article title='pq-Grams'>
  <author>Augsten</author>
  <author>Boehlen</author>
  <author>Gamper</author>
</article>
```
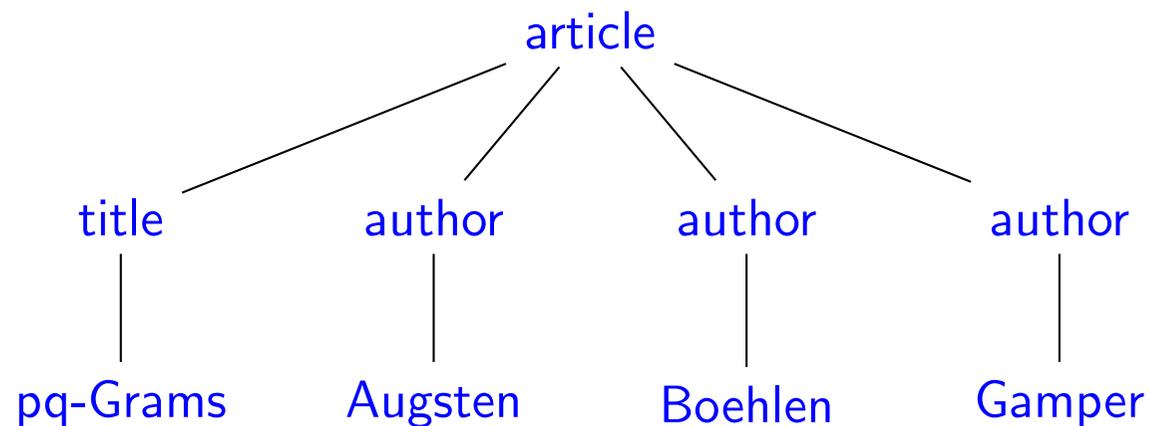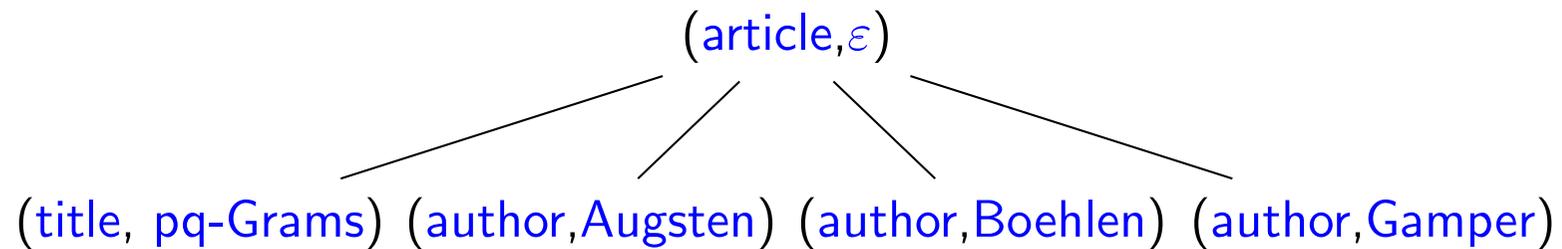
XML as a Double-Label Tree

- Node labels are pairs
- The XML document is encoded as a tree with:
  - XML element: node labeled with (tag-name,text-value)
  - XML attribute: node labeled with (attribute-name,text-value)
- Element nodes contain:
  - nodes of their sub-elements and attributes
- Attribute nodes are always leaves
- Element nodes without attributes or sub-elements are leaves
- Order:
  - sub-element nodes are ordered
  - attributes are not ordered (approach: see previous slide)
- Limitation: Can represent
  - *either* elements with sub-elements and/or attributes
  - *or* elements with a text value
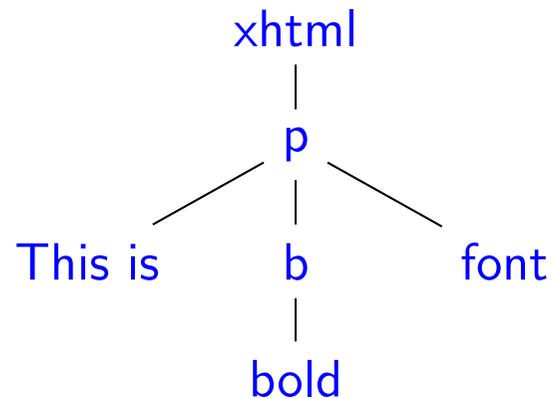
Example: XML as a Double-Label Tree

```
<article title='pq-Grams'>
  <author>Augsten</author>
  <author>Boehlen</author>
  <author>Gamper</author>
</article>
```
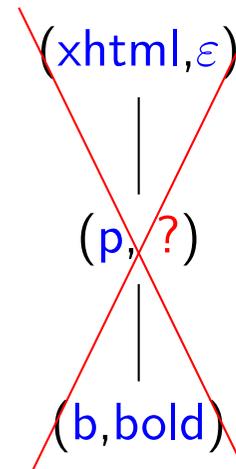
$(\text{article}, \varepsilon)$

$(\text{title}, \text{pq-Grams})$ $(\text{author}, \text{Augsten})$ $(\text{author}, \text{Boehlen})$ $(\text{author}, \text{Gamper})$

Example: Single- vs. Double-Label Tree

```
<xhtml>
  <p>This is <b>bold</b> font.</p>
</xhtml>
```

Single-Label Tree

xhtml

p

This is   b   font

bold

Double-Label Tree

(xhtml,ε)

(p, ?)

(b,bold)

Parsing XML

We discuss two popular parsers for XML:

- DOM – Document Object Model
- SAX – Simple API for XML
- StAX – Streaming API for XML

DOM – Document Object Model

- W3C[1] standard for accessing and manipulating XML documents
- Tree-based: represents an XML document as a tree (single-label tree with additional node info, e.g. node type)
- Elements, attributes, and text values are nodes
- DOM parsers load XML into main memory
  - random access by traversing tree :-)
  - large XML documents do not fit into main memory :-(

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[1]`http://www.w3schools.com/dom`

SAX / StAX

- SAX – Simple API for XML[2]
  - 'de facto" standard for parsing XML
- Event-based: reports parsing events (e.g., start and end of elements)
  - no random access :-(
  - you see only one element/attribute at a time
  - you can parse (arbitrarily) large XML documents :-)
- StAX - Streaming API for XML[3]
  - similar to SAX, but pull-based (vs. SAX: push)
  - pull: the client receives the next event on request
- Java API available for DOM, SAX, and StAX.
- For importing XML into a database: use SAX or StAX!

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[2]http://www.saxproject.org
[3]https://en.wikipedia.org/wiki/StAX