
 Exercise 1 - *Sorted Neighborhood Matching*.

2 Points

Consider the clean-clean entity resolution problem in Figure 1. Use the sorted neighborhood technique to produce *candidate pairs*, i.e., the record pairs that must be compared. Sort lexicographically in ascending order by *Name*, *PLZ*, and *YoB*, and use a window size of $w = 2$.

Figure 1: Clean-clean entity resolution problem.

ID	Name	ZIP	YoB	ID	Name	ZIP	YoB
a_1	Smith	5020	1993	b_1	Gruber	5034	1998
a_2	Christen	5020	1998	b_2	Smyth	5020	1993
a_3	Huber	5034	1993	b_3	Huber	5034	1949
a_4	Buber	5034	1949	b_4	Gruber	5020	2011
a_5	Gruber	5010	2011	b_5	Chirsten	5020	1998
				b_6	Huber	5034	1993

Name:

Matrikelnummer:

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Exercise 2 - *String Edit Distance*.

2 Points

Compute the edit distance between the strings `proof` and `porof`. Use the matrix produced by the dynamic programming algorithm to derive the shortest edit scripts and represent them with the gap representation.

Exercise 3 - q -Gram Distance.**2 Points**

Given the strings $x = \text{blabla}$ and $y = \text{alaba}$. Compute the q -gram distance and the normalized q -gram distance between x and y ($q = 2$).

Exercise 4 - Optimal Substructure of the String Edit Distance Problem.**2 Points**

Proof the optimal substructure property of the string edit distance problem:

Given a gap representation, $\text{gap}(x, y)$, between two strings x and y , such that the cost of $\text{gap}(x, y)$ is the string edit distance $\text{ed}(x, y)$. If we remove the last column of $\text{gap}(x, y)$, then the gap representation of the remaining columns, $\text{gap}(x', y')$, has cost $\text{ed}(x', y')$ between the resulting substrings, x' and y' .

Exercise 5 - Forest Distance Matrix.

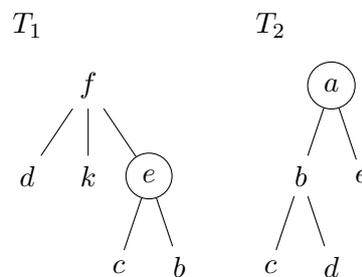
2 Points

Consider ordered trees T_1 and T_2 in Figure 2, forest distance matrix fd , and tree distance matrix td for the trees T_1 and T_2 .

- Compute the left-most leaf descendant arrays l_1, l_2 for trees T_1, T_2 , respectively.
- Fill the missing values for d_i and d_j into the forest distance matrix fd for the circled keyroot nodes in trees T_1 and T_2 .
- Circle the cell in the forest distance matrix that stores the distance between the prefixes $T_1[3..4]$ and $T_2[1..3]$.
- Cross all cells in the forest distance matrix and/or the tree distance matrix required to compute the distance between the prefixes $T_1[3..4]$ and $T_2[1..3]$.

fd	\vec{d}_j					
$d_i \downarrow$						

td	1	2	3	4	5
1					
2					
3					
4					
5					
6					

Figure 2: Two ordered trees T_1 and T_2 .

Exercise 6 - *Constrained Edit Distance*.

2 Points

Consider ordered trees T_1 and T_2 in Figure 3. Show an edit mapping between T_1 and T_2 that is *not* a valid constrained edit mapping. Explain.

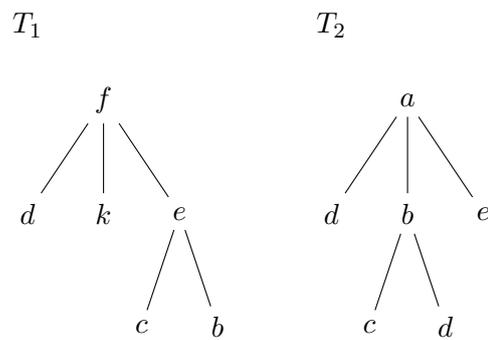


Figure 3: Two ordered trees T_1 and T_2 .

Exercise 7 - Binary Branch Distance and Lower Bound.**2 Points**

For the ordered trees T_1 and T_2 in Figure 4:

- Represent T_1 and T_2 as normalized binary trees and compute the binary branch distance.
- Based on the binary branch distance, what is the smallest value that the edit distance between T_1 and T_2 can adopt?

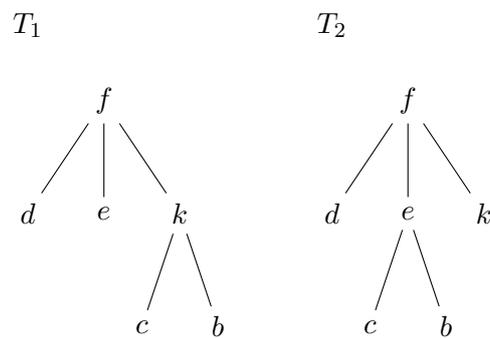


Figure 4: Two ordered trees T_1 and T_2 .

Exercise 8 - *Cosine Prefix Signature*.2 Points

Consider the collection $R = \{s_1, s_2, s_3, s_4\}$ of sets in Figure 5. Compute *prefix signatures* for all sets $s_i \in R$ for *cosine similarity* threshold $t = 0.9$.

Note: For the cosine similarity, $\text{Cos}(r, s)$, between two sets, r and s , the following holds:

$$\text{Cos}(r, s) \geq t \Rightarrow |r \cap s| \geq t^2|r|$$

$$\begin{aligned} s_1 &= \{A, C, B, D, F, E\} \\ s_2 &= \{D, E, F, B, A\} \\ s_3 &= \{B, D, F, G\} \\ s_4 &= \{C, B, G\} \end{aligned}$$

Figure 5: Set collection $R = \{s_1, s_2, s_3, s_4\}$.